





② EXPECTED  
# OF COMPARISONS =

$\{1, 2, 3, \dots, n\}$  SET TO  
 BE SORTED

$$= \sum_{i=1}^n \sum_{j=1}^{n-1} p_{ij} \in \Theta(n^2)$$

PROBABILITY  
 OF COMPARING

$i$  ~~VS~~  $j$   
~~VS~~  $b$

=

PROB. THAT IN  $[i, \dots, j]$   
 THE FIRST PIVOT WE CHOOSE IS  $i$   
 EITHER  $j$  OR  $j$

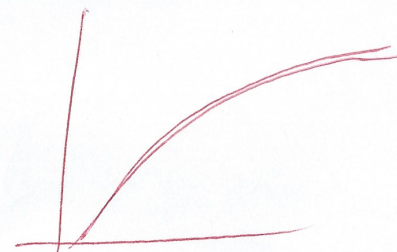
$j-i+1$  VALUES

$$\sum_{i=1}^n \sum_{j=1}^{n-1} \frac{2}{j-i+1}$$

~~$\sum_{k=1}^n \sum_{z=1}^{n-1} \frac{2}{z}$~~

HARMONIC SERIES

$$\frac{2}{j-i+1}$$



$$\sum_{z=1}^n \frac{1}{z} \approx \ln n + 0.566 \dots$$



# WHY NOT PREFERING QUICK SORT

• NOT STABLE

STABLE SORTING

DOS NOT EXCHANGE ELEMENT  
HAVING THE SAME VALUE

MERGE SORT → STABLE

FEW DISTINCT

• MANY ELEMENTS,  
FEW VALUES

$10^6$  ELEMENTS  
IN  $\{1, 2\}$

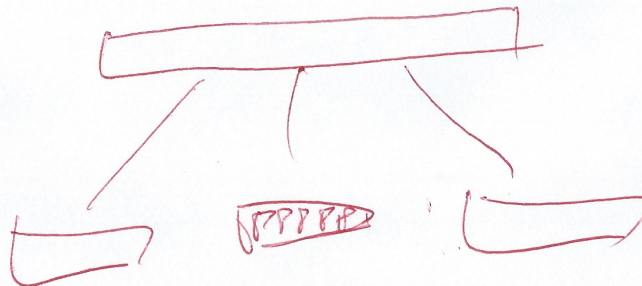
→  $O(n^2)$

THREE WAYS  
QUICK SORT

SMALLER THAN  
PIVOT

= PIVOT

LARGER THAN  
PIVOT

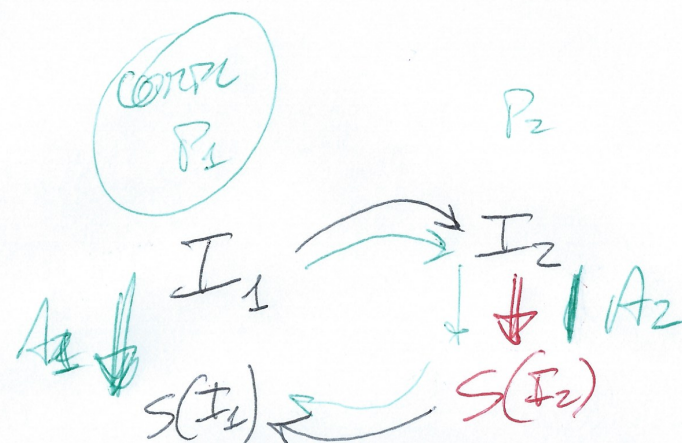
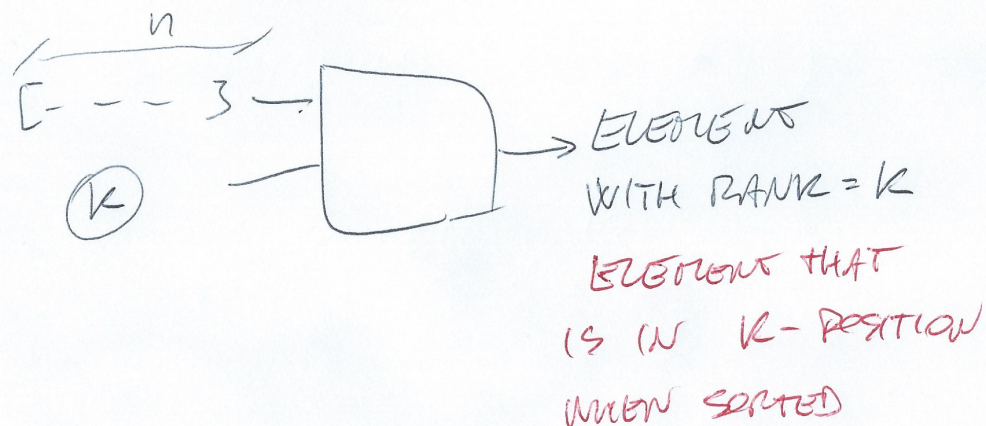




# ④ SELECTION PROBLEM (K-th ELEMENT PROBLEM)

- FIND MINIMUM (1-st ELEM)
- FIND MAXIMUM (n-th ELEM)
- FIND MEDIAN ( $\frac{n}{2}$ -th ELEM)

CAN DO BY SORTING  $\rightarrow O(n \log n)$



SELECTION PROBLEM  
NOT MORE DIFFICULT  
THAN SORTING

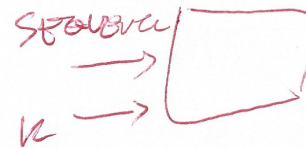
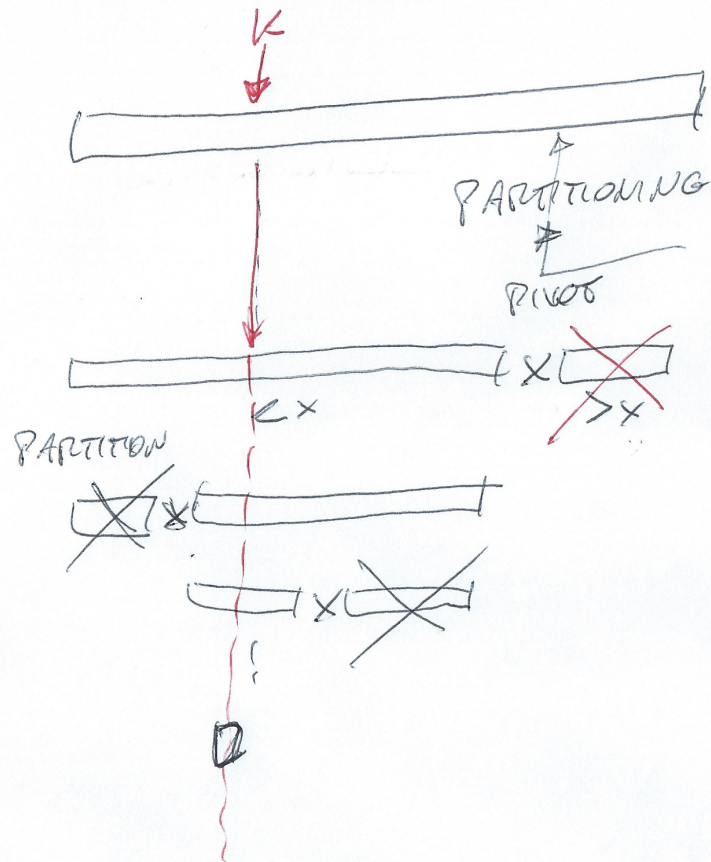
REDUCTION  
FROM  $P_1$   
TO  $P_2$

ALGORITHM FOR  $P_2$   
ALSO SOLVES  $P_1$



SELECTION K-th elem. IN  $\Theta(n)$  EXPECTED TIME

"PARTIAL" QUICK SORT



$START \leq K \leq END$

~~def sel~~

SELECTION (SEQ, START, END, K)

• BASE STEP {  $START = END$  }  $\rightarrow K$   
DONE

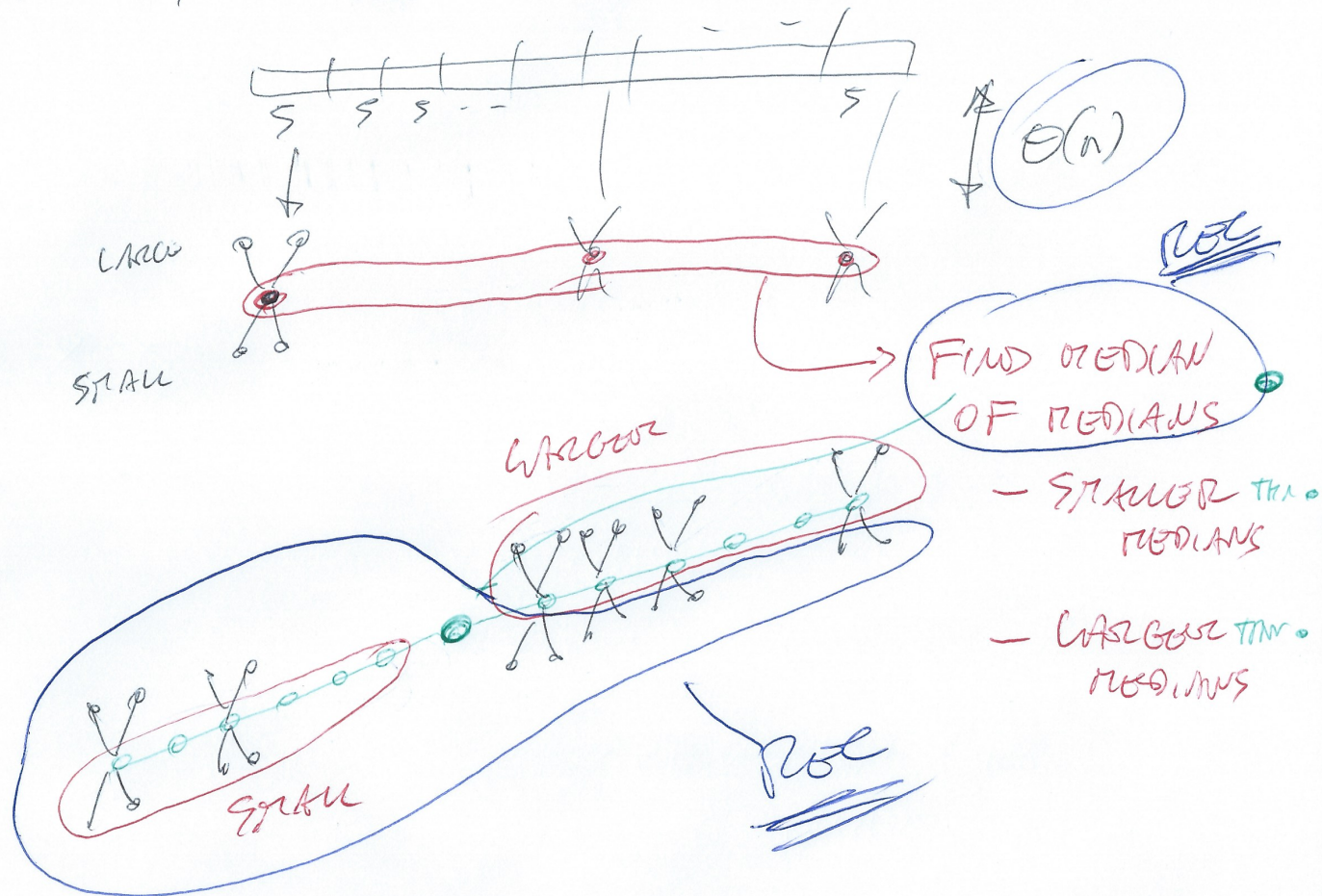
• PARTITION

$\rightarrow$  EITHER FOUND K-th

$\rightarrow$  RECURSE ON THE  
APPROPRIATE PORTION

6) Selection  $k$ -th elem.  $\Theta(n)$  Worst case time  
ALWAYS OPTIMAL

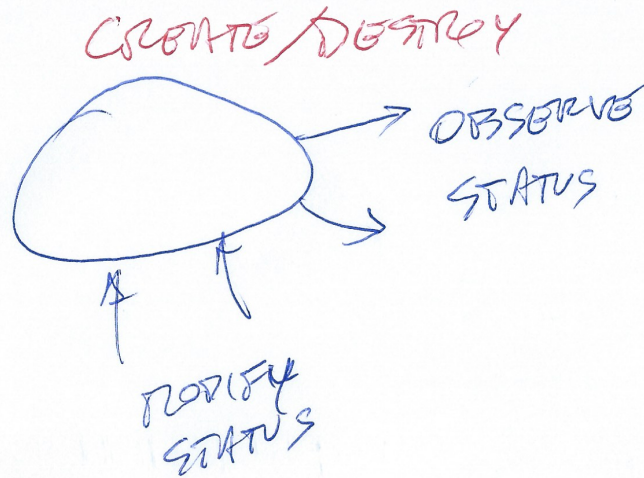
[BLUM, FLOYD, PRATT, RIVERA, TARJAN]





7)

ABSTRACT  
DATA  
TYPE (ADT)



## EXAMPLES

- ADD A VALUE TO
- ASK WHETHER  $x$  HAS BEEN EVER ADDED
- HOW MANY VALUES ARE THERE
- IS IT EMPTY?

SET

- CREATION

EMPTY SET

- ASK  $x \in S$

TRUE IFF OPERATION

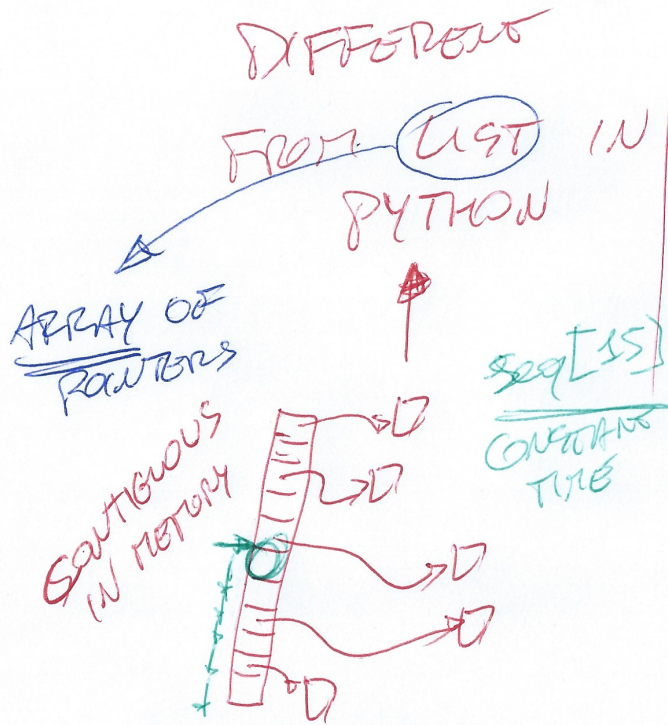
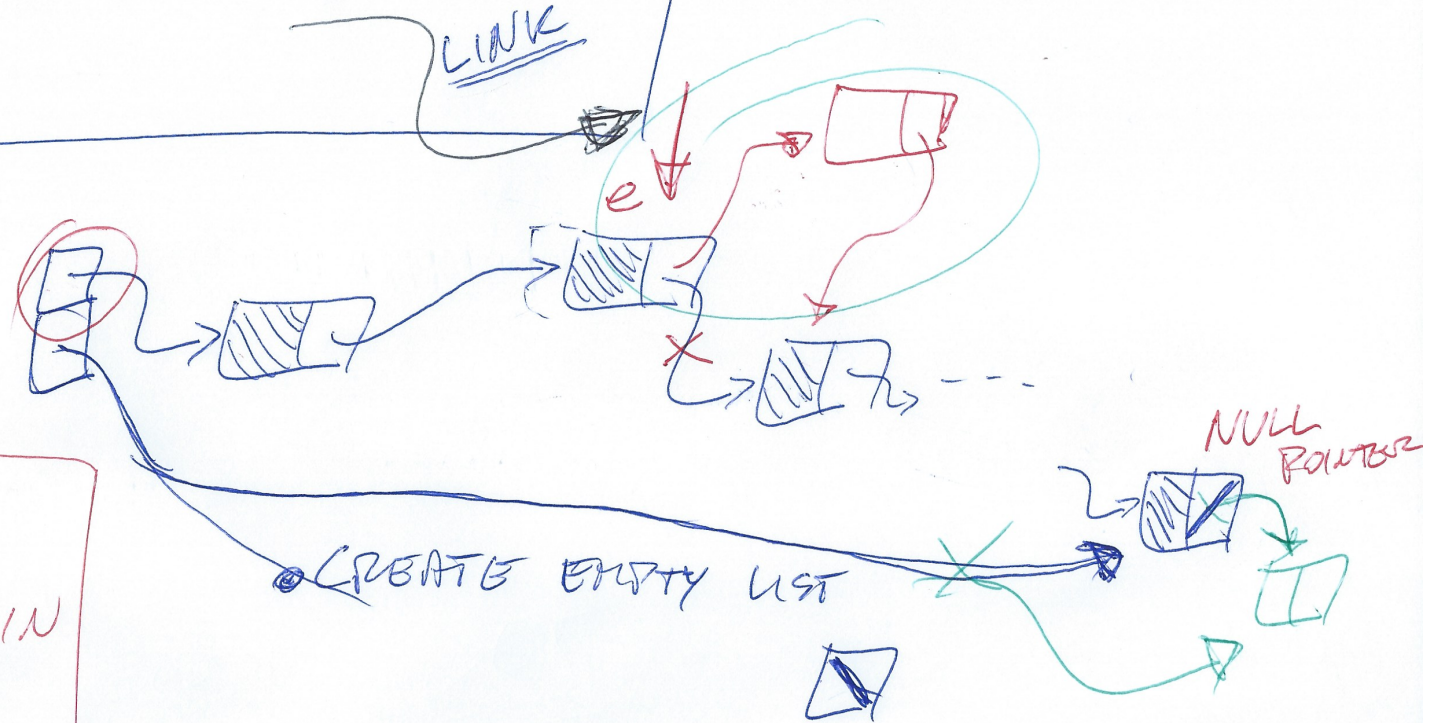
ADD( $x$ ) HAS BEEN PERFORMED BEFORE



⑧

# POINTER BASED (REFERENCES) STRUCTURES

## LINKED LIST



- CREATE EMPTY LIST
- GET FIRST ELEM. (IF EXISTS)  $O(1)$
- GET  $i$ -th ELEM. (if exists) SCAN LIST TO THE  $i$ -th.
- GET LENGTH.
- ADD ELEMENT 

TOP

AFTER e

END

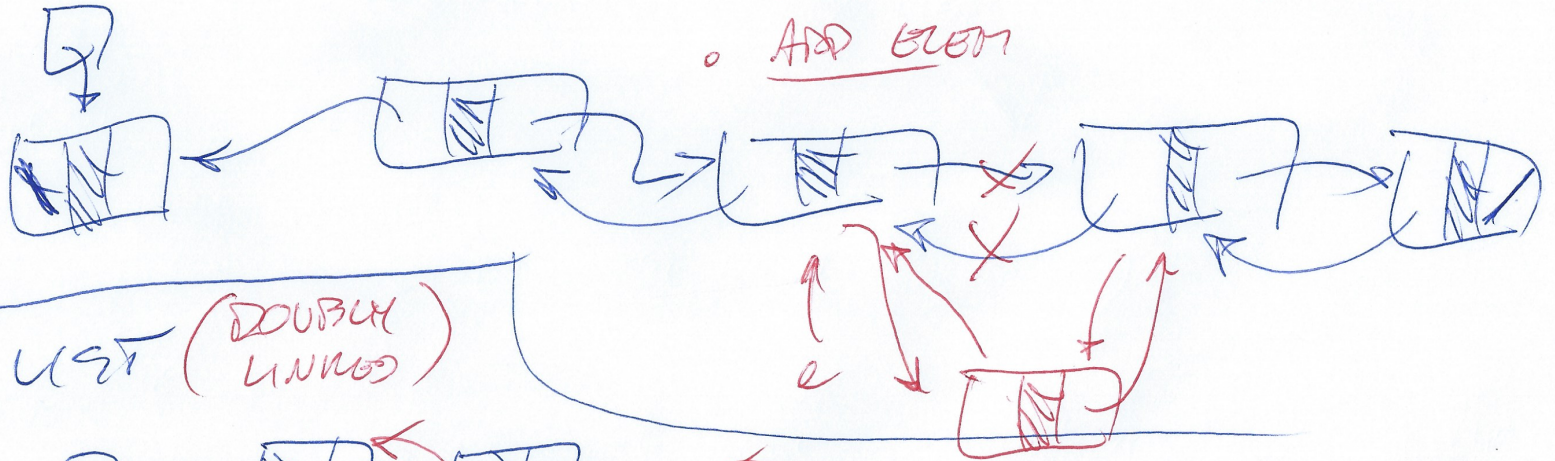


## LINKED LIST:

- ADD AFTER  $e \rightarrow \Theta(1)$  KNOWING A REF TO  $e$
- NEXT( $e$ )  $\Theta(1)$  KNOWING  $e$  REF
- PREV( $e$ ) ??? SCAN FROM THE START !!  
 $\Theta(n)$
- LAST

## DOUBLY LINKED LIST

- PREV( $e$ )  $\rightarrow \Theta(1)$
- ADD ELEM



## CIRCULAR LIST (DOUBLY LINKED)

